

# A Standard for Real-Time Smart Transducer Interface

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## Abstract

*In order to handle the inherent complexity of the multitude of available different transducer components, a generic interface approach is necessary. Such a universal smart transducer interface should provide precisely defined interfaces between smart transducers and their users, which are simple and precisely specified within the value domain and the temporal domain.*

*The Object Management Group adopted a smart transducer interface standard that incorporates (i) real-time characteristics and functionalities for the smart transducer network (ii) online diagnostic service capability (iii) support for start-up and dynamic configuration (iv) a uniform naming and addressing scheme for all relevant data in the smart transducer system (v) a generic interface that enables the smart transducer system to interact with other systems via CORBA.*

*This paper describes the main concepts and implementation experiences of this smart transducer interface. The approach integrates a time-triggered communication protocol with an appropriate access scheme, the so-called interface file system. This interface file system provides a unique naming and addressing scheme enabling access to internal transducer data via a CORBA gateway.*

## 1 Introduction

The decreasing cost of computing power enables decentralization of data processing by providing computational intelligence close to transducers (sensors and actuators). This enables data abstraction necessary for an abstract interface definition. Such an interface definition comprises a major element in software engineering for embedded systems avoiding an emerging crisis in embedded systems design due to the increasing number of different transducers in current embedded systems.

In embedded system applications with dramatically increasing numbers of transducers, as it is the case in the automotive market, which is marked by an ad-hoc

approach of system engineering with a multitude of different components (and different interface definitions), a generic approach can significantly reduce the development costs.

Some interface descriptions for embedded components exist already and are used extensively (e.g., CAN Kingdom [1], OSEK [2], and IEEE 1451.2 [3]), yet do address a precise interface description in the time and space domain and standardization is often supplier specific. System engineering of a CAN-based approach, e.g., can be inherently difficult and lead to unintended system behavior for certain CAN implementation (see e.g. Meschi et al. [4]).

In this paper we present a generic approach for the organization of transducers and communication between transducers that provides features in terms of real-time guarantees, complexity management, and maintainability. Smart transducers integrate an analog or digital sensor or actuator element and a local microcontroller that contains the interface circuitry, a processor, memory, and a network controller within a single unit. A smart transducer transforms the raw sensor signal to a standardized digital representation, checks and calibrates the signal, and transmits this digital signal via a standardized communication protocol to its users [5]. The possibility of reducing complexity using the concept of a smart transducer is of paramount benefit as basic element of a system engineering approach in embedded networks, which cannot be valued highly enough, though initially overlooked due to the current supply chain structure.

In addition to providing a building block for a successful management of a smooth software engineering process of embedded systems, a precise interface definition facilitates communication system design. We present a communication system with guaranteed timeliness and a deterministic description of communication traffic. Such a description of a distributed transducer system design can alleviate system development. Time partitioning mechanisms on a communication bus decrease the wiring loom complexity and provide additional benefits such as reduced weight, reduced wiring

costs, and improve the system reliability due to the decreased number of parts with high failure rates such as connectors.

In December 2000 the Object Management Group (OMG) called for a proposal (RFP) of a *smart transducer interface* (STI) that satisfies the following needs: (i) real-time characteristics and functionalities for the smart transducer network (ii) online diagnostic service capability (iii) support for start-up and dynamic configuration (iv) a uniform naming and addressing scheme for all relevant data in the smart transducer system (v) a generic interface that enables the smart transducer system to interact with other systems via a CORBA (Common Object Request Broker Architecture) gateway, and (vi) the support of communication interfaces available on current low-cost microcontrollers, e. g., UART ports. In response to this RFP, a time-triggered communication architecture with a well-defined interface to a CORBA environment has been submitted jointly by three companies with support of the Vienna University of Technology. This proposed STI standard has been adopted and published by the OMG as a world-wide standard in 2003 [6]. In this paper we describe this approach and present two case study implementations.

## 2 Related Work

A smart transducer interface should conform to a world-wide standard. Such a standard for real-time communication systems has been long sought, but vendors of existing were reluctant to support such a single common standard in fear of losing some of their competitive advantages.

In the field of fieldbus communication, several different fieldbus solutions have been developed and promoted. Some of these existing solutions have been combined and standardized. The European CENELEC standard was created by including all national standards into the standard parts EN 50170 (general purpose), EN 50254 (high efficiency), and EN 50325 (based on Controller Area Network (CAN)). In 1994, the two large fieldbus groups ISP (Interoperable Systems Project supported by Fisher-Rosemount, Siemens, Yokogawa, and others) and the WorldFIP (Flux Information Processus or Factory Instrumentation Protocol, supported by Honeywell, Bailey, and others) joined to form the Fieldbus Foundation (FF). It is the stated objective of the FF to develop a single interoperable fieldbus standard in cooperation with the International Electrotechnical Commission (IEC) and the Instrumentation Society of America (ISA). The IEC worked out the IEC 61158 standard, which is based on eight existing fieldbus solutions, among them Foundation Fieldbus, Profibus, and WorldFIP. The IEC and CENELEC standards have the great disadvantage that they still keep a diversity of very different solutions [7].

Standards for smart transducers have been developed apart from the fieldbus standardization efforts. The

IEEE 1451.2 standard [3] deals with the specification of interfaces for smart transducers. This standard defines electronic data sheets to describe the hardware interface and communication protocols of the smart transducer interface model. The IEEE 1451 standard includes an adequate naming/addressing scheme and supports the configuration of large transducer systems, but it lacks the explicit specification of real-time communication among the smart transducers.

## 3 Design Principles of STI

The architecture of the smart transducer interface has been guided by the following principles: (i) definition of common code spaces for naming, time and value domain, (ii) separation of interfaces according to their purpose, and (iii) support for integrating separately built and tested sub-systems.

### 3.1 Common Code Spaces for Naming, Time and Value Domain

On an abstract level, the purpose of a real-time smart transducer interface is the timely exchange of observations of real-time entities between the engaged subsystems across the provided interfaces. A real-time entity is a state variable of interest that has a name and a value at a particular instant. An observation [8] is thus an atomic triple:

$$\langle \text{name, observation instant, value} \rangle,$$

where *name* is an element of the common name space of real-time entities, the *observation* instant is a point in the time space and *value* is an element of the chosen value domain. An observation expresses that the referenced real-time entity possessed the stated value at the indicated instant. A synchronized global time among all nodes enables a common notion of time and is a prerequisite for meaningful timestamping of observations.

### 3.2 Interface Separation

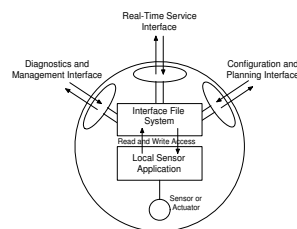


Fig. 1: The Three Interface Types to a Smart Transducer Node

If different user groups access the system for different purposes, they should only be provided with an interface to the information relevant for their respective purpose [9]. Therefore, interfaces for different purposes

may differ by the accessible information and in the temporal behavior of the access across the interface. As depicted in Figure 1 the STI specifies three different interface types to a smart transducer:

**DM interface:** This is a *diagnostic and management* interface. It establishes a connection to a particular smart transducer node and allows reading or modifying of specific IFS records. Most sensors need parametrization and calibration at startup and continuously collect diagnostic information to support the maintenance activities. For example a remote maintenance console can request diagnostic information from a certain sensor. The DM interface is usually not time-critical.

**CP interface:** The *configuration and planning* interface allows the integration and setup of newly connected nodes. It is used to generate the “glue” in the network that enables the components of the network to interact in the intended function. Usually, the CP interface is not time-critical.

**RS interface:** The *real-time service* interface performs a periodic communication with predictable timing behavior among the smart transducer nodes. Communicated data is usually data from sensors and for actuators. This view employs sensors for producing periodic observations of real-time entities in the environment. For example, a temperature sensor periodically sends the observed and locally preprocessed sensor value to the temporal firewall of the master. Since in TTP/A the time interval between sensing the environment and presenting the sensor value at the temporal firewall [10] of the master is known a priori, it is possible to perform a feed-forward state estimation of the sensor value at the sensor node in such a way, that the delivered sensor value is a good estimate of the real-time entity’s actual state at the point in time of delivery.

### 3.3 Temporal Composability

In many engineering disciplines, large systems are built by the constructive integration of well-specified and pre-tested subsystems, called components. During the system design phase, this requires a two-level design approach. At the overall system design level, the system integrator precisely defines the properties and interactions of the overall systems in the value and time domain. The cluster design serves as a requirements definition for the component design where local details can be defined independently.

The components are characterized by their physical parameters and the services they provide across well-specified interfaces. In a composable architecture, this integration should proceed without unintended side effects. To be composable, an architecture must adhere to four necessary principles with respect to the interfaces of nodes [11]:

**Independent Development of Nodes:** Nodes can only be designed independently of each other, if the architecture supports the exact specification of all node services provided to the system. The interface data structures must be precisely specified in the value domain and in the temporal domain and a proper conceptual interface model of the node service, as viewed by a user of the node, must be available.

**Stability of Prior Services:** The stability-of-prior-service principle ensures that the validated service of a node - both in the value domain and in the time domain - is not refuted by the integration of the node into a system.

**Performability of the Communication System:**

The performability-of-the-communication-system principle ensures that if  $n$  nodes are already integrated, the integration of the node  $n+1$  will not disturb the correct operation of the  $n$  already integrated nodes. A properly configured time-triggered communication system satisfies this requirement.

**Replica Determinism:** If fault tolerance is implemented by the replication of nodes, then the architecture and the nodes must support replica determinism. A set of replicated nodes is replica determinate if all the members of this set have the same externally visible state, and produce the same output messages at points in time that are at most an interval of  $d$  time units apart (as seen by an omniscient outside observer), where  $d$  is the duration necessary to replace a missing or erroneous message by a correct one.

## 4 Properties of the STI Standard

The STI standard defines a smart transducer system comprising several clusters with transducer nodes connected to a bus. Via a master node, each cluster is connected to a CORBA gateway. The master nodes of each cluster share a synchronized time that supports coordinated actions (e. g., synchronized measurements) over transducer nodes in several clusters. Each cluster can contain up to 250 smart transducers that communicate via a cluster-wide broadcast communication channel. There may be redundant shadow masters to support fault tolerance. One active master controls the communication within a cluster (in the following sections the term master refers to the active master unless stated otherwise). Since smart transducers are controlled by the master, they are called slave nodes. Figure 2 depicts an example for such a smart transducer system.

The interface file system is used to provide a unique addressing scheme for the interfaces. It is possible to monitor the smart transducer system via the CORBA interface without disturbing the real-time traffic.

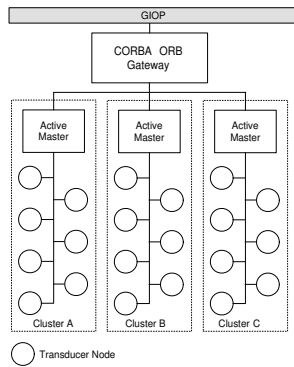


Fig. 2: Multi-Cluster Architecture with CORBA Gateway

The hardware requirements for smart transducer nodes are very flexible. The STI standard requires a minimum agreed set of services for a smart transducer implementation, thus supporting low-cost implementations of smart transducers, while allowing optional implementation of standard features. It is possible to fit a minimum STI implementation on an embedded micro-controller with 2k flash memory and 64 bytes of RAM memory.

#### 4.1 Interface File System

The information transfer between a smart transducer and its client is achieved by sharing information that is contained in an internal Interface File System (IFS), which is encapsulated in each smart transducer. The IFS provides a unique address scheme for transducer data, configuration data, self-describing information, and internal state reports of a smart transducer [5].

**Communication via temporal firewalls.** A time-triggered sensor bus will perform a periodical time-triggered communication to copy data from the IFS to the fieldbus and write received data into the IFS. Thus, the IFS is the source and sink for all communication activities. Furthermore, the IFS acts as a temporal firewall that decouples the local transducer application from the communication activities. A temporal firewall [10] is a fully specified interface for the unidirectional exchange of state information between a sender/receiver over a time-triggered communication system. The basic data and control transfer of a temporal firewall interface is depicted in Figure 3, showing the data and control flow between a sender and a receiver. The IFS at the sender forms the output firewall of the sender and the IFS of the receiver forms the input firewall of the receiver.

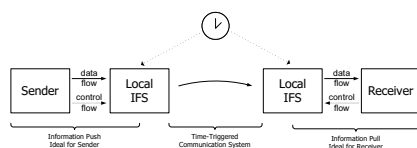


Fig. 3: Temporal Firewall

**Flow control using information push and pull paradigms.** The sender deposits its output information into its local IFS according to the information *push* paradigm, while the receiver must *pull* the input information out of its local IFS (non-consumable read) [12]. In the information push model the sender presses information on the receiver. It is ideal for the sender, because the sender can determine the instant for passing outgoing information to the communication system. The information pull model is ideal for the receiver, since tasks of the receiver will not be interrupted by incoming messages. The transport of the information is realized by a time-triggered communication system that derives its control signals autonomously from the progression of time. The instants when typed data structures are fetched from the senders IFS and the instant when this typed data structures are delivered to the receivers IFS are common knowledge of the sender and the receiver. A predefined communication schedule defines time, origin, and destination of each protocol communication. Thus, the IFS acts as a *temporally specified interface* that decouples the local transducer application from the communication task.

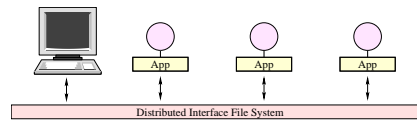


Fig. 4: Logical Network Structure

**Naming and addressing.** Each transducer can contain up to 64 files in its IFS. An IFS file is an index sequential array of up to 256 records. A record has a fixed length of four bytes (32 bits). An IFS record is the smallest addressable unit within a smart transducer system. Every record of an IFS file has a unique hierarchical address (which also serves as the global name of the record) consisting of the concatenation of the cluster name, the logical name, the file name, and the record name.

Besides access via the master node, the local applications in the smart transducer nodes are also able to execute a clusterwide application by communicating directly with each other. Figure 4 depicts the network view for such a clusterwide application.

The IFS of each smart transducer node can be accessed via the RS interface, the DM interface, and the CP interface for different purposes. All three interface types are mapped onto the fieldbus communication protocol, but with different semantics regarding timing and data protection.

#### 4.2 Fieldbus Communication Protocol

A time-triggered communication service following the specification of the STI has been implemented in the time-triggered fieldbus protocol TTP/A.

The bus allocation is done by a Time-Division Multiple-Access (TDMA) scheme. Communication is organized into rounds consisting of several TDMA slots. A slot is the unit for transmission of one byte of data. Data bytes are transmitted in a standard UART format. Each communication round is started by the master with a so-called fireworks byte. The fireworks byte defines the type of the round and is a reference signal for clock synchronization. The protocol supports eight different firework bytes encoded in a message of one byte using a redundant bit codesupporting error detection. Generally, there are two types of rounds:

**Multipartner round:** This round consists of a configuration dependent number of slots and an assigned sender node for each slot. The configuration of a round is defined in a datastructure called “RODL” (ROund Descriptor List). The RODL defines which node transmits in a certain slot, the operation in each individual slot, and the receiving nodes of a slot. RODLs must be configured in the slave nodes prior to the execution of the corresponding multipartner round. An example for a multipartner round is depicted in Figure 5.

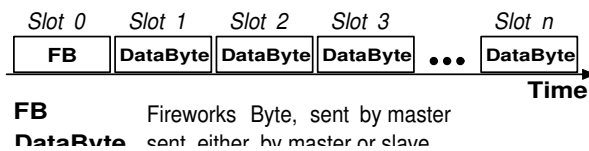


Fig. 5: Communication layout of a TTP/A multipartner round

**Master/slave round:** A master/slave round is a special round with a fixed layout that establishes a connection between the master and a particular slave for accessing data of the node’s IFS, e. g. the RODL information. In a master/slave round the master addresses a data record using a hierarchical IFS address and specifies an action like reading of, writing on, or executing that record.



Fig. 6: Recommended schedule with alternative multipartner and master/slave rounds

The master/slave rounds establish the DM and the CP interface to the transducer nodes. The RS interface is provided by periodical multipartner rounds. Master/slave rounds are scheduled periodically between multipartner rounds as depicted in Figure 6 in order to enable maintenance and monitoring activities during system operation without a probe effect [13].

## 5 Implementation Experiences

The following two case studies present a proof-of-concept of the STI standard. A configuration and maintenance tool presents additional means supporting convenient development and monitoring of applications.

### 5.1 Robot Arm

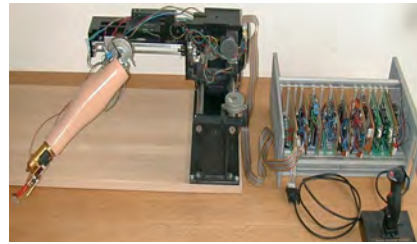


Fig. 7: Robot Arm

As a demonstrator for the STI we built a system with a robot arm. At the application level a human operator can control a prosthetic arm mounted on top of a linear thrust unit (see Figure 7). Simplicity of control for the user is established by the presence of intelligence in the demonstrator. Smart sensors yield information about the environmental conditions allowing avoidance of operating errors and obtaining precise control. Pressure sensors are present for determining the required grip force to grasp an object in order that no intervention of the human operator is needed to avoid slipping of an object. The robot arm is equipped with an angle sensor to allow limiting the opening of the elbow.

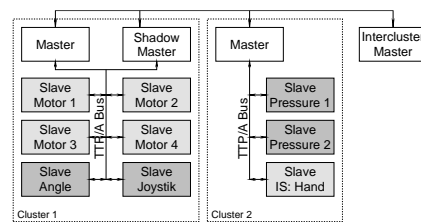


Fig. 8: Architecture of Robot Arm Transducer Network

The demonstrator was also built to investigate partitioning of nodes into distinct clusters.

The demonstrator consists of two clusters (see Figure 8). The first cluster contains the nodes for controlling the motors of the linear thrust units, the elbow, and the wrist. Nodes for retrieving the current angle of the elbow and the joystick commands are also placed in this cluster. A shadow master can take over control in case the primary master fails. Both masters are connected to the intercluster bus and act as intermediate systems. In addition to their TTP/A master role, they are slaves of a time-triggered backbone bus. The second cluster contains a node acting as an interface system for integrating the prosthetic hand into the demonstrator. Two nodes equipped with pressure sensors obtain measurements for grasping objects intelligently.

### 5.2 Autonomous Mobile Robot

Another implementation of the STI is shown by a model car that acts as an autonomous robot with sensory inputs. Figure 9 depicts the setup of this “smart car”.

The smart car contains a network of seventeen nodes whereof some of these nodes are implemented on very small microcontrollers to demonstrate the possibility of cheap slaves.

As indicated in Figure 10, the network comprises a smart car equipped with a suit of pivoted distance sensors, an electric drive and a steering unit. Distance sensors, servo motors for sensor pivoting, driving and steering units are all separate smart transducer nodes. Each node is implemented on a low-cost microcontroller and equipped with an STI.

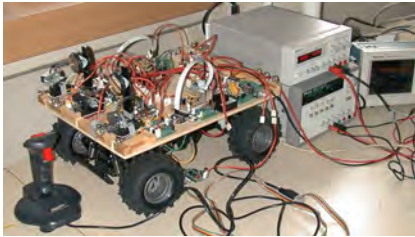


Fig. 9: Smart Car

The STI supports the integration of new connected smart transducer nodes into a system with predictable timing behavior. For example, it is possible to add extra “light” nodes to the car during operation.

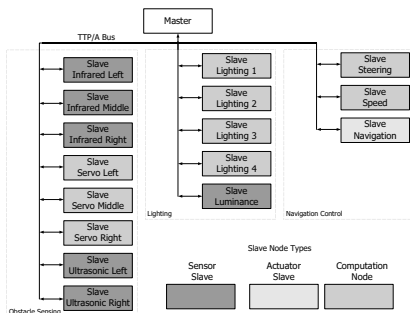


Fig. 10: Architecture of Smart Car Transducer Network

In order to achieve efficient operation two different operation modes are defined. The STI standard supports up to 6 user-definable independent communication modes, which support applications running different modes.

As long as no obstacles are detected within the sensors’ range, the car operates in “rabbit mode”. In this mode the car drives straight forward at full speed and two infrared sensors are aimed slightly outward the driving direction. The main detection of obstacles relies on two ultrasonic sensors. These are capable to report obstacles straight ahead of the car within a range of about 5m.

In case an obstacle is detected, the car switches to “turtle mode”. In this mode the car uses a communication schedule where all infrared sensors and pivot servos are serviced. The distance sensors are swivelled around by servo motors so that they are able to scan the area in front of the robot. The sensors generate a value that

corresponds to the distance of the object they are aimed at. The data stream provided by the distance sensors is taken over by a data processing node that fuses the perceptions from the distance sensors and the directions they are aimed at with a model of the robot environment. In this model the shapes of obstacles are stored and assigned with a probability value, that decreases with the progression of time and increases when the object is re-scanned. From this data a navigation node calculates the speed and the direction to provide this information to the speed and steering nodes.

### 5.3 Monitoring and Configuration Tool

Monitoring and debugging of distributed embedded real-time systems differ significantly from debugging and testing programs for desktop computers, because only few interfaces to the outside world are present. In addition, a distributed system contains multiple locations of control and therefore conventional breakpoint techniques result in an undesired modification of the timing behavior. This indeterministic effect is called the “Probe Effect” [13] or the “Heisenberg Uncertainty” [14] applied to software. Therefore, a vital property of a convenient monitoring environment is the absence of disturbances or intrusions on the system behavior. Users expect tools to avoid the probe effect and to incorporate a deterministic and reproducible behavior.

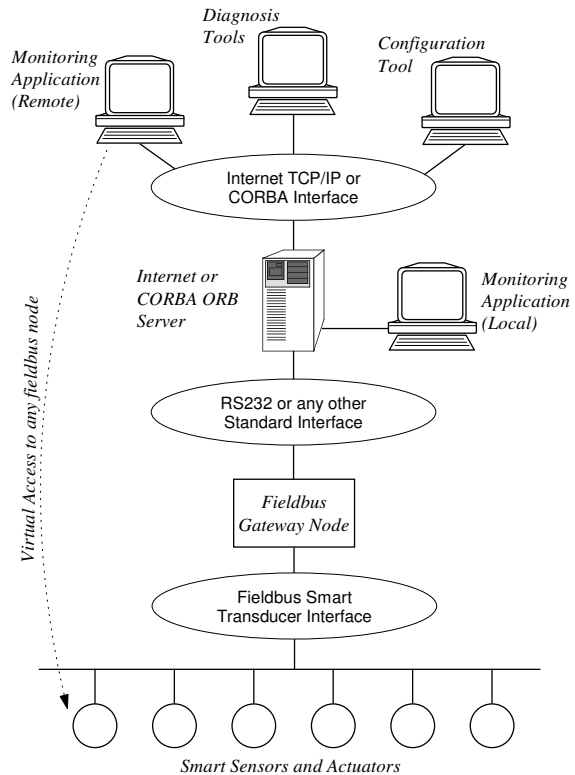


Fig. 11: System Architecture



We have developed a tool set for hard real-time environments, which establishes access via the DM and the CP interface. In contrast to existing tools, relevant node internal information is accessible without changing the behavior in either the temporal or value domain of the RT interface. This property allows evolutionary changes and monitoring activities of the system during real-time operation [15].

Figure 11 depicts the architecture of the monitoring approach. The fieldbus gateway node is connected to the fieldbus and contains the software and hardware to establish a connection via a standard PC interface, in our case an RS232 serial connection. Other possible connections are RS485, USB, IEEE 1394, or a wireless connection with IrDA or radio signals. The connection between fieldbus and server preserves real-time behavior. The server runs a CORBA communication module and supports local tools for monitoring and configuration. Because the timing behavior of the communication from the local tools into the fieldbus network provides a deterministic timing behavior, the local tools can support deterministic real-time monitoring. However, compared to access at fieldbus level data is accessible at a lower bandwidth.

The communication module running on the PC is a CORBA object. The CORBA middleware enables the transparent access of remote service applications to the fieldbus network through, e. g. the Internet. This interface is independent of the employed protocols and physical layers at fieldbus level.

## 6 Relation to ISO/OSI Model

Because on the timing requirements of control systems, fieldbus systems implement several layers of the OSI (Open Systems Interconnection) model at once, in most cases there is a representation for the OSI levels 1 (physical layer), 2 (data link layer), and 7 (application) whereas OSI layers 3...6 are usually void since they have no counterpart in the fieldbus world.<sup>1</sup>

The here presented STI standard defines OSI layers from application layer to link layer, whereas there are no features defined that would belong to layers 3...6. Figure 12 defines the relation between the OSI standard and the OMG smart transducer interface.

The IFS at the top level acts as transparent network interface for the application and as gateway to the CORBA network. The specified TDMA bus arbitration scheme and the codification of message length refer to the data link layer of the OSI model. According to the standard, the physical layer is open to various implementations, as long as the communication parameters of the chosen physical layer match the application's real-time requirements on timing and throughput. Thus, any compatible

<sup>1</sup> A notable exception is the LON (Local Operating Network) fieldbus, a communication system for building automation, which implements all 7 layers of the basic OSI model.

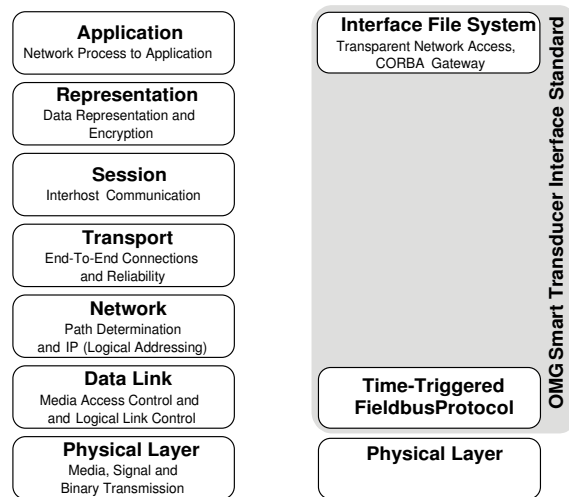


Fig. 12: Relation to layers of OSI model

physical layer standard such as RS 485, ISO 9141, CAN physical layer, IrDA, etc. may be used.

## 7 Discussion and Conclusion

One requirement stated in the request for proposal by the OMG has been the real-time capabilities of the smart transducer interface. The STI supports hard real-time communication by introducing a time-triggered communication scheme that is a priori specified before the RS interface of the system comes into operation. Generally, time-triggered systems require an increased effort in the design phase of the system, but provide an easier verification of the temporal correctness. Since time-triggered systems are designed according to the principle of resource adequacy [16], it is guaranteed that sufficient computing resources are available to handle the specified peak load scenario. On the other hand, time-triggered systems are often blamed for their bad flexibility. The STI overcomes this limitation by introducing means to configure the interaction of the components via the CP interface.

The RS interface provides composability, guaranteed timeliness, and hides the components' internals. The DM and CP interfaces involve inherently event-triggered activities, which require an event-triggered communication service. These interfaces cannot invalidate the temporal behavior of the RS interface and support full access to component internals – as required by a maintenance engineer.

The specification of interfaces should be complete and of minimal cognitive complexity. Cognitive complexity can be minimized by restricting interactions via carefully designed interfaces and by providing access restrictions. The kind of information that must be available via an interface depends on the purpose of the particular interface. For example, a properly designed operational interface hides component internals, thereby

allowing a component to form a meaningful abstraction. The corresponding operational interface specification stipulated during architecture design should incorporate a precise specification of a component's inputs and outputs in both the temporal and the value domain. A maintenance engineer on the other hand, might require access to intermediate computational results for locating the origin of an incorrect system behavior.

The STI standard meets the requirement for complete interfaces of minimal cognitive complexity by introducing three different types of interfaces of a component. The separation (RS, DM, and CP interface) is done according to the interface purpose, the necessary level of visibility of component internal information, and the type of the temporal interaction patterns. Such a separation minimizes complexity in contrast to a single interface type incorporating support for all possible interactions.

The STI standard also specifies the provision of the three interface types through a CORBA server. However, currently there is no CORBA architecture for effectively supporting the temporal requirements to establish the RS interface. Current priority-based approaches like real-time CORBA require complete knowledge about all other service requests and their corresponding priority values in the whole CORBA network when guarantees about the temporal behavior are required. Furthermore, the availability of a global notion of time allows to record the instant of the acquisition of a real-time entity's state in each observation.

As a proof of concept, we have implemented two case studies of the STI. The first case study comprises a demonstrator with a robot arm that is instrumented by a smart transducer network partitioned into two clusters. The second case study is an autonomous mobile robot, that shows the integration of new nodes and efficient communication despite of static communication schedules. The case studies show that the STI standard is an interesting option for a wide range of networked sensing and control applications. The STI provides many features that are required by fieldbus applications for automotive or automation industries. Supported features are the real-time capability, the encapsulation of the node's internals, and a universal address space with the interface file system. The STI can be implemented on low-cost Commercial-off-the-Shelf (COTS) hardware and supports various bus media types.

## 8 Acknowledgments

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